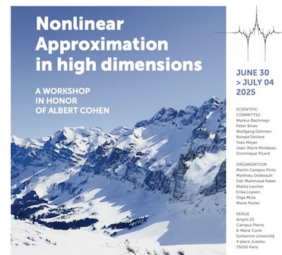


# How many (noisy) measurements are needed to learn a vector?

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# The problem

Assume that you want to **recover an arbitrary vector**  $x \in \mathbb{R}^m$ , up to some error  $\varepsilon > 0$  in some norm  $\|\cdot\|$ , where  $m \in \mathbb{N}$  can be large.

You know nothing about  $x$ , you can only compute certain measurements  $\lambda_1(x), \dots, \lambda_n(x) \in \mathbb{R}$  that you can choose.

**How many measurements do you need?**

(We do not assume a bound on a norm of  $x$  yet.)

# Linear information

**Less than  $m$  linear measurements are useless.**

Proof:

Let  $N = (\lambda_1, \dots, \lambda_n)$  with  $n < m$ , where the  $\lambda_k$  are linear functionals.

Then, for  $y = N(x)$ , there is a whole affine subspace  $V$  of  $\mathbb{R}^m$  with  $\dim(V) \geq m - n$  such that  $N(v) = y$  for all  $v \in V$ .

So no matter how you choose your approximation  $\hat{x} = \Phi(y) \in \mathbb{R}^m$ , you may be arbitrarily far away from the true value of  $x$ .

# Continuous information

The same is true for any continuous measurement map  $N: \mathbb{R}^m \rightarrow \mathbb{R}^n$ .  
This follows from the Borsuk-Ulam theorem:

Borsuk-Ulam theorem (~1930)

For any  $N \in C(\mathbb{R}^m, \mathbb{R}^n)$  with  $n < m$  and  $R > 0$ , there are  $x$  and  $\tilde{x}$  at distance  $2R$  such that  $N(x) = N(\tilde{x})$ .

That is, you cannot distinguish between these vectors and hence cannot guarantee an error less than  $R$ .

So, again, **less than  $m$  continuous measurements are useless.**

# Adaptive information

But why should we fix the measurement maps  $\lambda_k$  in advance?

It is probably better to use the already obtained information  $\lambda_1(x), \dots, \lambda_{k-1}(x)$  and choose  $\lambda_k$  based on it.

We call such measurements **adaptive**.

(Precisely, we allow  $\lambda_k(x) = \lambda_k(x; \lambda_1(x), \dots, \lambda_{k-1}(x), \lambda_1, \dots, \lambda_{k-1})$ .)

# Adaptive linear information

“Unfortunately”, the fact remains that

**less than  $m$  linear measurements are useless,  
also if we choose them adaptively.**

Proof:

As before, for  $y = N(x)$ , there is a whole affine subspace  $V \subset \mathbb{R}^m$  such that  $\lambda_j(v) = y_j$  for all  $j \leq n < m$  and  $v \in V$ .

Thus, for any  $v \in V$ , you would have chosen the same functionals  $\lambda_j$  and obtained the same information  $y = N(v)$ .

Hence, you cannot distinguish between all the elements of  $V$ .

# Adaptive continuous information

The remaining case:

**Can we achieve something with less than  $m$   
adaptively chosen continuous measurements?**

History: When we started to work on this problem with David Krieg and Erich Novak during the summer 2024, we studied the “simplest example”  $x \in \{x \in \mathbb{R}^3 : \|x\|_2 \leq 1\}$ , i.e.,  $m = 3$ , and  $n = 2$ .

Starting with the functional  $\ell_1(x) = x_1 - |x_3|$ , one can improve over all nonadaptive methods (which have worst case error 1).

# Surprisingly, $\log(m)$ measurements are enough

## Theorem

[Krieg/Novak/U '25]

Let  $m \in \mathbb{N}$  and  $\varepsilon > 0$ . The algorithm  $R_m^\varepsilon$  described below uses at most  $n(m) := \lceil \log_2(m) \rceil + 1$  adaptive, 1-Lipschitz-continuous measurements and satisfies for all  $x \in \mathbb{R}^m$  that

$$\|x - R_m^\varepsilon(x)\| \leq \varepsilon.$$

This implies

$$\|R_m^\varepsilon(x) - R_m^\varepsilon(y)\| \leq \|x - y\| + 2\varepsilon,$$

which might be considered as some kind of *stability*.

Again: There cannot be a continuous algorithm! (Borsuk-Ulam)



# The algorithm

The algorithm is based on a **coloring** of  $\mathbb{R}^m$  in the following way:

Consider a partition

$$\mathbb{R}^m = \bigcup_{i \in \mathbb{N}} D_i,$$

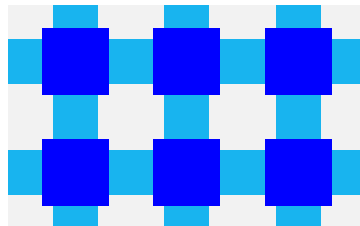
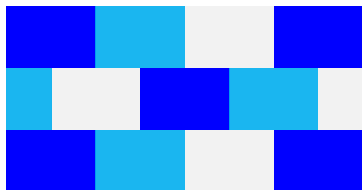
a (coloring) map  $t: \mathbb{N} \rightarrow \{1, 2, \dots, m+1\}$ , and some  $c > 0$  with:

- ①  $\text{diam}(D_i) \leq 1$  for each  $i$ ;
- ② if  $i \neq j$  and  $t(i) = t(j)$  then  $\text{dist}(D_i, D_j) \geq c$ ,

where diameter ( $\text{diam}$ ) and distance ( $\text{dist}$ ) are with respect to the given norm on  $\mathbb{R}^m$ .

# Colorings of $\mathbb{R}^m$ with separated colors

The following illustrations show two colorings of the plane with three colors. The second is easily generalized to higher dimensions.



Less than  $m + 1$  colors do not work.

This is related to the **Nagata dimension** of  $\mathbb{R}^m$ .

# The algorithm: First stage and bisection

To find an  $\varepsilon$ -approximation of  $x \in \mathbb{R}^m$  it is enough to find an index  $i^*$  such that  $x$  is in  $\varepsilon D_{i^*}$ . Define by

$$I_r := \{i \in \mathbb{N} : t(i) = r\} \quad \text{and} \quad E_r = \bigcup_{j \in I_r} \varepsilon D_j$$

the **set of points with color  $r$**  in the  $\varepsilon$ -scaled partition.

A continuous measurement of the form  $\lambda_J(x) = \text{dist}\left(x, \bigcup_{j \in J} E_j\right)$  with  $J \subset \{1, \dots, m+1\}$ , tells us whether  $x$  **has any of the colors in  $J$** .

We use  $n = \lceil \log_2(m+1) \rceil = n(m) - 1$  such functionals and **bisection to find a color of  $x$** , i.e.,  $r^* = t(i^*)$  with  $x \in \overline{E}_{r^*}$ .

## Algorithm: Second stage

Now we can determine a correct index  $i^*$  with  $x \in \varepsilon \overline{D}_{i^*}$  using any continuous functional  $\lambda^*$  for which the images  $\lambda^*(\varepsilon \overline{D}_i)$  for  $i \in I_{r^*}$  are pairwise disjoint. An example is given by

$$\lambda^*(x) = \max_{i \in I_{r^*}} \left\{ \frac{c\varepsilon}{2i} - \text{dist}(x, \varepsilon D_i) \right\}.$$

By construction,  $x$  is in the closure of  $\varepsilon D_{i^*}$  with  $i^* = \frac{c\varepsilon}{2\lambda^*(x)}$ .

The output  $R_m^\varepsilon(x)$  of the algorithm can be any element from  $\varepsilon D_{i^*}$ .

# Discussion

**Did we really construct a clever algorithm  
for the recovery of  $x \in \mathbb{R}^m$  if  $m$  is large?**

We made two basic assumptions:

- 1 If  $\lambda: \mathbb{R}^m \rightarrow \mathbb{R}$  is 1-Lipschitz, then we can compute values  $\lambda(x)$ , i.e., all Lipschitz functionals are admissible as information.
- 2 The information can be chosen adaptively, i.e.,  $\lambda_{k+1}$  may depend on the (already computed) values  $y_i = \lambda_i(x)$  for  $i = 1, 2, \dots, k$ .

Adaption is widespread and can be easily implemented on a computer.  
The use of arbitrary Lipschitz measurements is more problematic.

# Power of adaption for continuous information

We want to **approximate**  $f \in F$  in the norm of  $Y$ , where  $F \subset Y$  is subset of (general) metric space  $Y$ . We allow **algorithms**  $A_n: F \rightarrow Y$  of the form

$$A_n(f) = \Phi(\lambda_1(f), \dots, \lambda_n(f))$$

with adaptively chosen continuous measurements  $\lambda_i: F \rightarrow \mathbb{R}$  and an arbitrary reconstruction map  $\Phi: \mathbb{R}^n \rightarrow Y$ .

We consider the **minimal worst-case error** of algorithms that use at most  $n$  continuous measurements, i.e.,

$$e_n^{\text{cont}}(F) := \inf_{A_n} \sup_{f \in F} d_Y(f, A_n(f)),$$

where the infimum ranges over all adaptive  $A_n$  as above.

# Power of adaption for continuous information II

We compare with the **manifold widths** of  $F$  (in  $Y$ ), i.e.,

$$\delta_n(F) := \inf_{\substack{N \in C(F, \mathbb{R}^n) \\ \Phi \in C(\mathbb{R}^n, Y)}} \sup_{f \in F} d_Y(f, \Phi(N(f))).$$

$\leadsto$  minimal errors of **non-adaptive “continuous algorithms”**

## Theorem

[KNU '25]

Let  $F \subset Y$  be a subset of a metric space  $Y$  and  $n \geq 2$ . Then,

$$e_n^{\text{cont}}(F) \leq \delta_{2^{n-2}}(F).$$

# Noisy information

**What if measurements are only known  
up to some noise  $\delta > 0$ ?**

Precisely,  $(y_1, \dots, y_n) \in \mathbb{R}^n$  is the information about  $f \in F$  which satisfies

$$\left| y_i - \lambda_i(f; y_1, \dots, y_{i-1}) \right| \leq \delta, \quad 1 \leq i \leq n,$$

for some functionals  $\lambda_i(\cdot; y_1, \dots, y_{i-1}): F \rightarrow [-1, 1]$  and  $\delta < 1$ .

(Interpretation: *ideal* measurements  $\lambda_i$ ; machine precision  $\delta$ ; normalization  $[-1, 1]$ )

We denote by  $e_n^?(F, \delta)$  with  $? \in \{\text{lin}, \text{cont}, \text{arb}\}$  the minimal error of **arbitrary algorithms using ? measurements that are noisy.**



# Linear noisy information

There is lot of work and books on **noisy linear measurements**.

Here, we only state

$$e_n^{\text{lin}}(F, \delta) \geq \delta \quad \text{for all } n \in \mathbb{N}.$$

That is, for fixed  $\delta > 0$ ,

**there cannot be a (possibly non-linear) algorithm  
based on noisy linear measurements  
with error going to zero.**

# Continuous noisy information

For noisy continuous measurements  $\lambda_i: F \rightarrow [-1, 1]$ , we obtain an upper bound with the **entropy numbers** of  $F \subset Y$ :

$$\varepsilon_n(F) := \inf\{\varepsilon > 0: F \text{ can be covered by } 2^n \text{ balls (in } Y) \text{ of radius } \varepsilon\}.$$

## Theorem

[Krieg/Novak/Plaskota/U '25]

Let  $F \subset Y$  be a subset of a metric space  $Y$ ,  $n \in \mathbb{N}$  and  $\delta < 1$ .

Then,

$$e_n^{\text{cont}}(F, \delta) \leq \varepsilon_n(F).$$

- upper bound is independent of  $\delta < 1$  (!?)
- $e_n^{\text{cont}}(F, \delta) = \varepsilon_0(F)$  for  $\delta \geq 1$
- We did not find a “good” lower bound for  $e_n^{\text{cont}}(F, \delta)$ .

# Arbitrary noisy information

We consider **arbitrary (possibly non-continuous) measurements**.

Here, the minimal errors are even **characterized by the  $\varepsilon_n$** :

## Theorem

[Krieg/Novak/Plaskota/U '25]

Let  $F \subset Y$  be a subset of a metric space  $Y$ ,  $n \in \mathbb{N}$  and  $\delta < 1$ .

Then,

$$\varepsilon_{n(k_\delta+1)}(F) \leq e_n^{\text{arb}}(F, \delta) \leq \varepsilon_{nk_\delta}(F)$$

with

$$k_\delta := \lceil \log_2(1/\delta + 1) \rceil - 1.$$

↪ gain of nonlinear adaptive measurements is limited for fixed  $\delta > 0$

# An example

We illustrate these results by the example of approximating vectors from the unit ball  $F = B_p^m$  of  $\ell_p^m$  in the norm of  $Y = \ell_q^m$ .

We know

$$\varepsilon_n(B_p^m, \ell_q^m) \asymp \left( \frac{\log(m/n + 1)}{n} \right)^{\frac{1}{p} - \frac{1}{q}}$$

for  $\log(m) \leq n \leq m$  and  $p \leq q$ .

We will only discuss  $p = 2$  and  $q = \infty$ .

# An example II

We denote the number of measurements for  $\ell_\infty$ -error  $\varepsilon \in (0, 1)$  by

$$n^?(\varepsilon, \delta) := \inf \left\{ n : e_n^?(B_2^m, \delta) \leq \varepsilon \right\}.$$

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We denote the number of measurements for  $\ell_\infty$ -error  $\varepsilon \in (0, 1)$  by

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We have that

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- $n^{\text{cont}}(\varepsilon, 0) \leq \lceil \log_2(m+1) \rceil$  for all  $\varepsilon > 0$ ,



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- $n^{\text{lin}}(\varepsilon, \delta) \gtrsim n^{\text{lin}}(\varepsilon, 0) \gtrsim m$  for  $\varepsilon > \delta$ ,
- $n^{\text{cont}}(\varepsilon, 0) \leq \lceil \log_2(m+1) \rceil$  for all  $\varepsilon > 0$ ,
- $\log_{1/\delta}(m) \cdot \varepsilon^{-2} \lesssim n^{\text{cont}}(\varepsilon, \delta) \lesssim \log_2(m) \cdot \varepsilon^{-2}$

# Final remarks

- For  $Y = \ell_q^m$ ,  $q \in \{1, \infty\}$ , the functionals are **piecewise linear**.
- Neural networks...
- For  $\ell_2$ -approximation in  $F = B_1^m$ , we have

$$n^{\text{lin}}(\varepsilon, \delta) \approx n^{\text{cont}}(\varepsilon, \delta) \quad \text{for } \varepsilon < \delta.$$

A particularly interesting question:

## Open problem

What about **other classes of measurements?**

(just norms; convex, homogeneous, smooth)

# Thank you, Albert!



# Thank you, Albert!



# Lipschitz noisy information

Let us consider **Lipschitz-continuous measurements**.

(The result depends on the Lipschitz constant.)

## Theorem

[Krieg/Novak/Plaskota/U '25]

For any  $\delta \cdot \varepsilon_0(F) < 1$ , we have

$$\frac{\delta}{L} \leq e_n^{\text{Lip}}(F, \delta) \leq \varepsilon_n(F) + \frac{n\delta}{L \cdot \varepsilon_0(F)},$$

where  $e_n^{\text{Lip}}$  is the  $n$ -th minimal error with  $L$ -Lipschitz functionals.

In particular,

$$e_n^{\text{cont}}(F, \delta) \leq \varepsilon_n(F).$$